1. Modular Arithmetic.

1. Modular Arithmetic. Clock Math!!!

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- 2. Inverses for Modular Arithmetic: Greatest Common Divisor.

- 1. Modular Arithmetic. Clock Math!!!
- 2. Inverses for Modular Arithmetic: Greatest Common Divisor. Division!!!

- 1. Modular Arithmetic. Clock Math!!!
- 2. Inverses for Modular Arithmetic: Greatest Common Divisor. Division!!!
- 3. Euclid's GCD Algorithm.

- 1. Modular Arithmetic. Clock Math!!!
- 2. Inverses for Modular Arithmetic: Greatest Common Divisor. Division!!!
- 3. Euclid's GCD Algorithm. A little tricky here!

Modular Arithmetic.

Applications: cryptography, error correction.

Theorem: If $d|x$ and $d|y$, then $d|(y-x)$.

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 $x = ad, y = bd,$

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Proof: $x = ad$, $y = bd$, $(x - y) = (ad - bd) = d(a - b) \implies d|(x - y).$

П

Theorem: If $d|x$ and $d|y$, then $d|(y-x)$.

Proof: $x = ad$, $y = bd$, $(x - y) = (ad - bd) = d(a - b) \implies d|(x - y).$

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П

Theorem: If $d|x$ and $d|y$, then $d|(y-x)$.

Proof: $x = ad$, $y = bd$, $(x - y) = (ad - bd) = d(a - b) \implies d|(x - y).$

Theorem: Every number *n* ≥ 2 can be represented as a product of primes.

Proof: Either prime, or $n = a \times b$, and use strong induction. (Uniqueness? Later.)

П

Г

Poll

What did we use in our proofs of key ideas?

- (A) Distributive Property of multiplication over addition.
- (B) Euler's formula.
- (C) The definition of a prime number.
- (D) Euclid's Lemma.

Poll

What did we use in our proofs of key ideas?

- (A) Distributive Property of multiplication over addition.
- (B) Euler's formula.
- (C) The definition of a prime number.
- (D) Euclid's Lemma.
- (A) and (C)

Next Up.

Modular Arithmetic.

If it is 1:00 now.

If it is 1:00 now. What time is it in 2 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00. $101 = 12 \times 8 + 5.$

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

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 $101 = 12 \times 8 + 5.$

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Clock time equivalent up to addition of any integer multiple of 12.

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Custom is only to use the representative in $\{12,1,\ldots,11\}$

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What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system. Clock time equivalent up to addition of any integer multiple of 12.

Custom is only to use the representative in $\{12,1,\ldots,11\}$ (Almost remainder, except for 12 and 0 are equivalent.)
This is Thursday is September 15, 2022.

This is Thursday is September 15, 2022. What day is it a year from now?

This is Thursday is September 15, 2022. What day is it a year from now? on September 15, 2022?

This is Thursday is September 15, 2022. What day is it a year from now? on September 15, 2022? Number days.

This is Thursday is September 15, 2022. What day is it a year from now? on September 15, 2022? Number days.

0 for Sunday, 1 for Monday, . . . , 6 for Saturday.

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Today: day 4.

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Today: day 4.

5 days from then.

This is Thursday is September 15, 2022. What day is it a year from now? on September 15, 2022? Number days. 0 for Sunday, 1 for Monday, . . . , 6 for Saturday.

Today: day 4.

5 days from then. day 9

This is Thursday is September 15, 2022. What day is it a year from now? on September 15, 2022? Number days. 0 for Sunday, 1 for Monday, . . . , 6 for Saturday.

Today: day 4.

5 days from then. day 9 or day 2

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Today: day 4.

5 days from then. day 9 or day 2 or Tuesday.

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Today: day 4.

5 days from then. day 9 or day 2 or Tuesday.

25 days from then.

This is Thursday is September 15, 2022. What day is it a year from now? on September 15, 2022? Number days. 0 for Sunday, 1 for Monday, . . . , 6 for Saturday.

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5 days from then. day 9 or day 2 or Tuesday.

25 days from then. day 29

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5 days from then. day 9 or day 2 or Tuesday. 25 days from then. day 29 or day 1. $29 = (7)4 + 1$

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two days are equivalent up to addition/subtraction of multiple of 7.

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11 days from then is day 1 which is Monday!

This is Thursday is September 15, 2022. What day is it a year from now? on September 15, 2022? Number days. 0 for Sunday, 1 for Monday, . . . , 6 for Saturday. Today: day 4. 5 days from then. day 9 or day 2 or Tuesday. 25 days from then. day 29 or day 1. $29 = (7)4 + 1$ two days are equivalent up to addition/subtraction of multiple of 7. 11 days from then is day 1 which is Monday!

What day is it a year from then?

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What day is it a year from then? Next year is not a leap year.

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What day is it a year from then?

Next year is not a leap year. So 365 days from then.

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What day is it a year from then? Next year is not a leap year. So 365 days from then. Day 4+365 or day 369.

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What day is it a year from then? Next year is not a leap year. So 365 days from then. Day 4+365 or day 369. Smallest representation:

subtract 7 until smaller than 7.

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What day is it a year from then? Next year is not a leap year. So 365 days from then. Day 4+365 or day 369. Smallest representation:

subtract 7 until smaller than 7. divide and get remainder.

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5 days from then. day 9 or day 2 or Tuesday. 25 days from then. day 29 or day 1. $29 = (7)4 + 1$ two days are equivalent up to addition/subtraction of multiple of 7. 11 days from then is day 1 which is Monday!

What day is it a year from then? Next year is not a leap year. So 365 days from then. Day 4+365 or day 369. Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

369/7

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What day is it a year from then?

Next year is not a leap year. So 365 days from then.

Day 4+365 or day 369.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

369/7 leaves quotient of 52 and remainder 5.

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What day is it a year from then?

Next year is not a leap year. So 365 days from then.

Day 4+365 or day 369.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

369/7 leaves quotient of 52 and remainder 5. $369 = 7(52) + 5$

This is Thursday is September 15, 2022.

What day is it a year from now? on September 15, 2022? Number days.

0 for Sunday, 1 for Monday, . . . , 6 for Saturday.

Today: day 4.

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What day is it a year from then?

Next year is not a leap year. So 365 days from then.

Day 4+365 or day 369.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

369/7 leaves quotient of 52 and remainder 5. $369 = 7(52) + 5$ or September 15, 2022 is a Friday.

This is Thursday is September 15, 2022.

What day is it a year from now? on September 15, 2022? Number days.

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Today: day 4.

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Next year is not a leap year. So 365 days from then.

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Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

369/7 leaves quotient of 52 and remainder 5. $369 = 7(52) + 5$ or September 15, 2022 is a Friday.

80 years?

80 years? 20 leap years.

80 years? 20 leap years. 366×20 days

80 years? 20 leap years. 366×20 days 60 regular years.

80 years? 20 leap years. 366×20 days 60 regular years. 365×60 days
80 years? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4.

80 years? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4+366 \times 20+365 \times 60$.

80 years? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4+366\times20+365\times60$. Equivalent to?

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Hmm.

What is remainder of 366 when dividing by 7?

80 years? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4+366\times20+365\times60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

80 years? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4+366\times20+365\times60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7?

80 years? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4+366\times20+365\times60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

80 years? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4+366\times20+365\times60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

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Years and years...
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What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
 What is remainder of 365 when dividing by 7? 1
Today is day 4.
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Years and years...
```
Hmm.

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
 What is remainder of 365 when dividing by 7? 1
Today is day 4.
```
Get Day: $4+2\times20+1\times60$

```
Years and years...
```
Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 4.

Get Day: $4 + 2 \times 20 + 1 \times 60 = 104$

```
Years and years...
```
Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 4. Get Day: $4 + 2 \times 20 + 1 \times 60 = 104$

Remainder when dividing by 7?

```
Years and years...
```

```
80 years? 20 leap years. 366 \times 20 days
 60 regular years. 365 \times 60 days
Today is day 4.
It is day 4+366\times20+365\times60. Equivalent to?
```

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
 What is remainder of 365 when dividing by 7? 1
Today is day 4.
  Get Day: 4 + 2 \times 20 + 1 \times 60 = 104Remainder when dividing by 7? 104 = 14 \times 7
```

```
Years and years...
```

```
80 years? 20 leap years. 366 \times 20 days
 60 regular years. 365 \times 60 days
Today is day 4.
It is day 4+366\times20+365\times60. Equivalent to?
```

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What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
 What is remainder of 365 when dividing by 7? 1
Today is day 4.
  Get Day: 4 + 2 \times 20 + 1 \times 60 = 104Remainder when dividing by 7? 104 = 14 \times 7 + 6.
```

```
Years and years...
```

```
80 years? 20 leap years. 366 \times 20 days
 60 regular years. 365 \times 60 days
Today is day 4.
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```

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
 What is remainder of 365 when dividing by 7? 1
Today is day 4.
  Get Day: 4 + 2 \times 20 + 1 \times 60 = 104Remainder when dividing by 7? 104 = 14 \times 7 + 6.
  Or September 15, 2102 is Saturday!
```

```
Years and years...
```

```
80 years? 20 leap years. 366 \times 20 days
 60 regular years. 365 \times 60 days
Today is day 4.
It is day 4+366\times20+365\times60. Equivalent to?
```

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What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
 What is remainder of 365 when dividing by 7? 1
Today is day 4.
  Get Day: 4 + 2 \times 20 + 1 \times 60 = 104Remainder when dividing by 7? 104 = 14 \times 7 + 6.
  Or September 15, 2102 is Saturday!
```
Further Simplify Calculation:

```
Years and years...
```

```
80 years? 20 leap years. 366 \times 20 days
 60 regular years. 365 \times 60 days
Today is day 4.
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 What is remainder of 365 when dividing by 7? 1
Today is day 4.
  Get Day: 4 + 2 \times 20 + 1 \times 60 = 104Remainder when dividing by 7? 104 = 14 \times 7 + 6.
  Or September 15, 2102 is Saturday!
```
Further Simplify Calculation:

20 has remainder 6 when divided by 7.

```
Years and years...
```

```
80 years? 20 leap years. 366 \times 20 days
 60 regular years. 365 \times 60 days
Today is day 4.
It is day 4+366\times20+365\times60. Equivalent to?
```

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What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
 What is remainder of 365 when dividing by 7? 1
Today is day 4.
  Get Day: 4 + 2 \times 20 + 1 \times 60 = 104Remainder when dividing by 7? 104 = 14 \times 7 + 6.
  Or September 15, 2102 is Saturday!
```
Further Simplify Calculation:

20 has remainder 6 when divided by 7.

60 has remainder 4 when divided by 7.

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Years and years...
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Today is day 4.
  Get Day: 4 + 2 \times 20 + 1 \times 60 = 104Remainder when dividing by 7? 104 = 14 \times 7 + 6.
  Or September 15, 2102 is Saturday!
```
Further Simplify Calculation:

20 has remainder 6 when divided by 7. 60 has remainder 4 when divided by 7. Get Day: $4 + 2 \times 6 + 1 \times 4 = 20$.

```
Years and years...
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  Or September 15, 2102 is Saturday!
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Further Simplify Calculation:

```
20 has remainder 6 when divided by 7.
  60 has remainder 4 when divided by 7.
Get Day: 4 + 2 \times 6 + 1 \times 4 = 20.
 Or Day 6.
```

```
Years and years...
```

```
80 years? 20 leap years. 366 \times 20 days
 60 regular years. 365 \times 60 days
Today is day 4.
It is day 4+366\times20+365\times60. Equivalent to?
```

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
 What is remainder of 365 when dividing by 7? 1
Today is day 4.
  Get Day: 4 + 2 \times 20 + 1 \times 60 = 104Remainder when dividing by 7? 104 = 14 \times 7 + 6.
  Or September 15, 2102 is Saturday!
```

```
Further Simplify Calculation:
```
20 has remainder 6 when divided by 7.

60 has remainder 4 when divided by 7.

Get Day: $4 + 2 \times 6 + 1 \times 4 = 20$.

Or Day 6. September 15, 2102 is Saturday.

```
Years and years...
```

```
80 years? 20 leap years. 366 \times 20 days
 60 regular years. 365 \times 60 days
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"Reduce" at any time in calculation!

x **is congruent to** *y* **modulo** *m* or " $x \equiv y \pmod{m}$ " if and only if $(x - y)$ is divisible by *m*.

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 $\{ \ldots, -7, 0, 7, 14, \ldots \}$

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Can calculate with representative in $\{0,\ldots,m-1\}$.

x (mod m) or mod (x, m)

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2x = 3 \implies (\frac{1}{2}) \cdot 2x = (\frac{1}{2}) \cdot 3 \implies x = \frac{3}{2}.
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Can solve $4x = 5 \pmod{7}$. $2.4x = 2.5$ (mod 7)

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Poll

Mark true statements.

(A) Mutliplicative inverse of 2 mod 5 is 3 mod 5.

- (B) The multiplicative inverse of $((n-1)$ (mod *n*) = $((n-1)$ (mod *n*)).
- (C) Multiplicative inverse of 2 mod 5 is 0.5.
- (D) Multiplicative inverse of $4 = -1 \pmod{5}$.
- (E) (−1)*x*(−1) = 1. Woohoo.

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(C) is false. 0.5 has no meaning in arithmetic modulo 5.

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Very different for elements with inverses.

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Bijection

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All the images are distinct. \implies unique pre-image for any image.

 $x = 2, m = 4.$

If $gcd(x,m) = 1$. Then the function $f(a) = xa$ mod *m* is a bijection. One to one: there is a unique pre-image(single *x* where $y = f(x)$.) Onto: the sizes of the domain and co-domain are the same. $x = 3, m = 4.$ $f(1) = 3(1) = 3$ (mod 4), $f(2) = 6 = 2 \pmod{4}$, $f(3) = 1 \pmod{3}$. Oh yeah. $f(0) = 0$ (mod 3).

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Not a bijection.

Poll

Which is bijection?

(A) $f(x) = x$ for domain and range being $\mathbb R$ (B) $f(x) = ax \pmod{n}$ for $x \in \{0, ..., n-1\}$ and $gcd(a, n) = 2$ (C) $f(x) = ax \pmod{n}$ for $x \in \{0, ..., n-1\}$ and $gcd(a, n) = 1$

Poll

Which is bijection?

(A) $f(x) = x$ for domain and range being $\mathbb R$ (B) $f(x) = ax \pmod{n}$ for $x \in \{0, ..., n-1\}$ and $gcd(a, n) = 2$ (C) $f(x) = ax \pmod{n}$ for $x \in \{0, ..., n-1\}$ and $gcd(a, n) = 1$ (B) is not.

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Inverses

Next up.
Next up.

Next up. Euclid's Algorithm.

Next up.

Euclid's Algorithm. Runtime.

Next up.

Euclid's Algorithm. Runtime. Euclid's Extended Algorithm.

Does 2 have an inverse mod 8?

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20 / 43

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Fact: If *d*|*x* and *d*|*y* then *d*|(*x* + *y*) and *d*|(*x* - *y*).

Notation: *d*|*x* means "*d* divides *x*" or $x = kd$ for some integer k . **Fact:** If *d*|*x* and *d*|*y* then *d*|(*x* + *y*) and *d*|(*x* - *y*). Is it a fact?

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More divisibility

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 (i f (= y 0)x
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Excursion: Value and Size.

Before discussing running time of gcd procedure...

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Before discussing running time of gcd procedure... What is the value of 1,000,000?

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Assume $log_2 1,000,000$ is 20 to the nearest integer. **Mark what's true.**

Poll.

Assume $log_2 1,000,000$ is 20 to the nearest integer. **Mark what's true.**

- (A) The size of 1,000,000 is 20 bits.
- (B) The size of 1,000,000 is one million.
- (C) The value of 1,000,000 is one million.
- (D) The value of 1,000,000 is 20.

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(A) and (C).

Poll

Which are correct?

(A)
$$
gcd(700,568) = gcd(568,132)
$$

(B) $gcd(8,3) = gcd(3,2)$
(C) $gcd(8,3) = 1$
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(The second is less than the first.)

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After $2\log_2 x = O(n)$ recursive calls, argument x is 1 bit number.

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Didn't necessarily need to do gcd. Runtime proof still works.

Finding an inverse?

We showed how to efficiently tell if there is an inverse.

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We showed how to efficiently tell if there is an inverse. Extend euclid to find inverse.

Euclid's GCD algorithm.

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Computes the gcd(*x*,*y*) in *O*(*n*) divisions.

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Computes the gcd(*x*,*y*) in *O*(*n*) divisions.

For *x* and *m*, if $gcd(x, m) = 1$ then *x* has an inverse modulo *m*.

Multiplicative Inverse.

GCD algorithm used to tell **if** there is a multiplicative inverse.

Multiplicative Inverse.

GCD algorithm used to tell **if** there is a multiplicative inverse. How do we **find** a multiplicative inverse?

Modular Arithmetic: $x \equiv y \pmod{N}$ if $x = y + kN$ for some integer *k*.

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 $(3)12 + (-1)35 = 1.$

 $a = 3$ and $b = -1$. The multiplicative inverse of 12 (mod 35) is 3.

gcd(35,12)

```
gcd(35,12)
 gcd(12, 11) ;; gcd(12, 35%12)
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How did gcd get 11 from 35 and 12?

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Algorithm finally returns 1.
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Get 1 from 12 and 11.

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ext-qcd(x,y)if y = 0 then return(x, 1, 0)
     else
         (d, a, b) := ext-qcd(y, mod(x, y))return (d, b, a - floor(x/y) * b)
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Claim: Returns (d, a, b) : $d = gcd(a, b)$ and $d = ax + by$. Example:

ext-gcd(35,12)

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Claim: Returns (d, a, b) : $d = gcd(a, b)$ and $d = ax + by$. Example: $a - |x/y| \cdot b =$

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ext-gcd(35,12)
  ext-gcd(12, 11)
    ext-qcd(11, 1)ext-qcd(1,0)return (1,1,0) ;; 1 = (1)1 + (0)0
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```
Claim: Returns (d, a, b): d = gcd(a, b) and d = ax + by.
Example: a - |x/y| \cdot b = 1 - |11/1| \cdot 0 = 1
```

```
ext-gcd(35,12)
 ext-gcd(12, 11)
   ext-qcd(11, 1)ext-qcd(1,0)return (1,1,0) ;; 1 = (1)1 + (0)0return (1,0,1) ;; 1 = (0)11 + (1)1
```

```
ext-qcd(x, y)if y = 0 then return(x, 1, 0)
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Claim: Returns (d, a, b): d = gcd(a, b) and d = ax + by.
Example: a− ⌊x/y⌋ ·b =
0− ⌊12/11⌋ ·1 = −1
```

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```

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Claim: Returns (d, a, b) : $d = gcd(a, b)$ and $d = ax + by$. Example: *a*− ⌊*x*/*y*⌋ ·*b* = 1− ⌊35/12⌋·(−1) = 3

```
ext-qcd(35,12)ext-gcd(12, 11)
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Theorem: Returns (d, a, b) , where $d = gcd(a, b)$ and

 $d = ax + by$.

Proof: Strong Induction.¹

¹ Assume *d* is $gcd(x, y)$ by previous proof.

Proof: Strong Induction.¹ **Base:** ext-gcd(*x*, 0) returns ($d = x, 1, 0$) with $x = (1)x + (0)y$.

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 $\mathsf{Recursively:}\ \ d = ay + b(x - \lfloor \frac{x}{y} \rfloor \cdot y) \implies d = bx - (a - \lfloor \frac{x}{y} \rfloor b)y$ Returns (*d*,*b*,(*a*− <u>∣</u> *y*^{*y*} *j* · *b*)).

Hand Calculation Method for Inverses.

Example: $gcd(7, 60) = 1$.

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7(1) + 60(0) = 7
$$

$$
7(-8) + 60(1) = 4
$$

$$
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$$

\n
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\n
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\n
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\n
$$
7(-17) + 60(2) = 1
$$

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Confirm:

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Confirm: $-119 + 120 = 1$

Conclusion: Can find multiplicative inverses in *O*(*n*) time!

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Internet Security.

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Internet Security. Public Key Cryptography: 512 digits.

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Internet Security: Next Week.